

# ECS455 Chapter 2

## Cellular Systems

### 2.4 Traffic Handling Capacity<sup>''</sup> and Erlang B Formula

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# Capacity Concept: A Revisit

- Q: If I have  $m$  channels per cell, is it true that my cell can support only  $m$  user?
- A: Yes and No
- Let's try one example.
- How often do you make a call?
  - 3 calls a day, on average. ←  $\lambda$  user
- How long is the call?
  - 10 mins (per call), on average. ←  $H = \frac{1}{\mu}$
- So, one person uses

$$3 \frac{\text{calls}}{\text{day}} \times \frac{10 \text{ mins}}{\text{call}} = \frac{30 \text{ mins}}{\text{day}} = \frac{30 \text{ mins}}{24 \times 60 \text{ mins}} = \frac{1}{48} \text{ [Erlang]}$$



# Capacity Concept: A Revisit

- If we can “give” the time that “User 1” is idle to other users,
  - then one channel can support 48 users!!  
(48x “capacity”)
- True? (Not quite)
- 48 users is possible if we have a way to manipulate all 48 users to not make calls when another user is using the channel.
- Real users access the channel randomly.  
(Call initiation/request times are random.)
- If we allow  $>1$  users, then we (the users) will have to deal with congestion.

# New Concepts

- Using  $m$  as the capacity of a cell is too small.
- We can let more than one user share a channel by using it at different times.
- The number of users that a cell can support can then exceed  $m$ .
- Call initiation times are random
- **Blocked calls**
- Probability of (call) blocking  $P_b$ 
  - the likelihood that a call is blocked because there is no available channel.
  - 1%, 2%, 5%

# Trunking

- Allow a large number ( $n$ ) of users to **share** the relatively small number of channels in a cell (or a sector) by providing access to each user, **on demand**, from a **pool** of available channels.
- Exploit the **statistical behavior** of users.
- Each user is allocated a channel on a per call basis, and upon termination of the call, the previously occupied channel is immediately returned to the pool of available channels.

# Common Terms (1)

- **Traffic Intensity**: Measure of channel time utilization (traffic load / amount of traffic), which is the average channel occupancy measured in **Erlangs**. In our example,
  - Dimensionless
  - Denoted by  $A$ . one user utilizes  $A_u = \frac{1}{48}$  Erlang
- **Holding Time**: Average duration of a typical call. If we have  $n=10$  users in the pool, then they utilize
  - Denoted by  $H = 1/\mu$ . = 10 mins  $A = \frac{10}{48}$  Erlang.
- **Request Rate**: The average number of call requests per unit time. Denoted by  $\lambda$ .  $\lambda_u = 3$  calls/day  
 $\lambda = 10 \times 3 = 30$  calls/day.
- Use  $A_u$  and  $\lambda_u$  to denote the corresponding quantities for one user.
- Note that  $A = nA_u$  and  $\lambda = n\lambda_u$  where  $n$  is the number of users supported by the pool (trunked channels) under consideration.

# Common Terms (2)

- **Blocked Call:** Call which cannot be completed at time of request, due to congestion. Also referred to as a **lost call**.
- **Grade of Service (GOS):** A measure of congestion which is specified as the probability of a call being blocked (for Erlang B).  
 $P_b$   $P_b \leq 0.02$
- The **AMPS** cellular system is designed for a GOS of 2% blocking. This implies that the channel allocations for cell sites are designed so that 2 out of 100 calls will be blocked due to channel occupancy during the busiest hour.

# M/M/m/m Assumption

- **Blocked calls cleared**
  - Offers **no queuing** for call requests.
  - For every user who requests service, it is assumed there is **no setup time** and the user is given immediate access to a channel if one is available.
  - If **no channel** are available, the requesting user is **blocked** without access and is **free to try again later**.
- **Calls arrive as determined by a *Poisson process***.
- There are **memoryless arrivals** of requests, implying that all users, including blocked users, may request a channel at any time.
- There are an **infinite number of users** (with finite overall request rate).
  - The finite user results always predict a smaller likelihood of blocking. So, assuming infinite number of users provides a conservative estimate.
- **The duration of the time that a user occupies a channel is exponentially distributed**, so that longer calls are less likely to occur.
- There are **m channels** available in the trunking pool.
  - For us,  $m =$  the number of channels for a cell ~~(S)~~ or for a sector



# Erlang B Formula

$$P_b = \frac{A^m}{m! \sum_{i=0}^m \frac{A^i}{i!}}$$

$m$  = Number of trunked channels

Call blocking probability

$A$  = traffic intensity or load [Erlangs]

$$= \frac{\lambda}{\mu}$$

$\lambda$  = Average # call attempts/requests per unit time

$$= \lambda \times \frac{1}{\mu}$$

$$= \lambda \times H$$

$\frac{1}{\mu} = H$  = Average call length

*not a build-in function.*

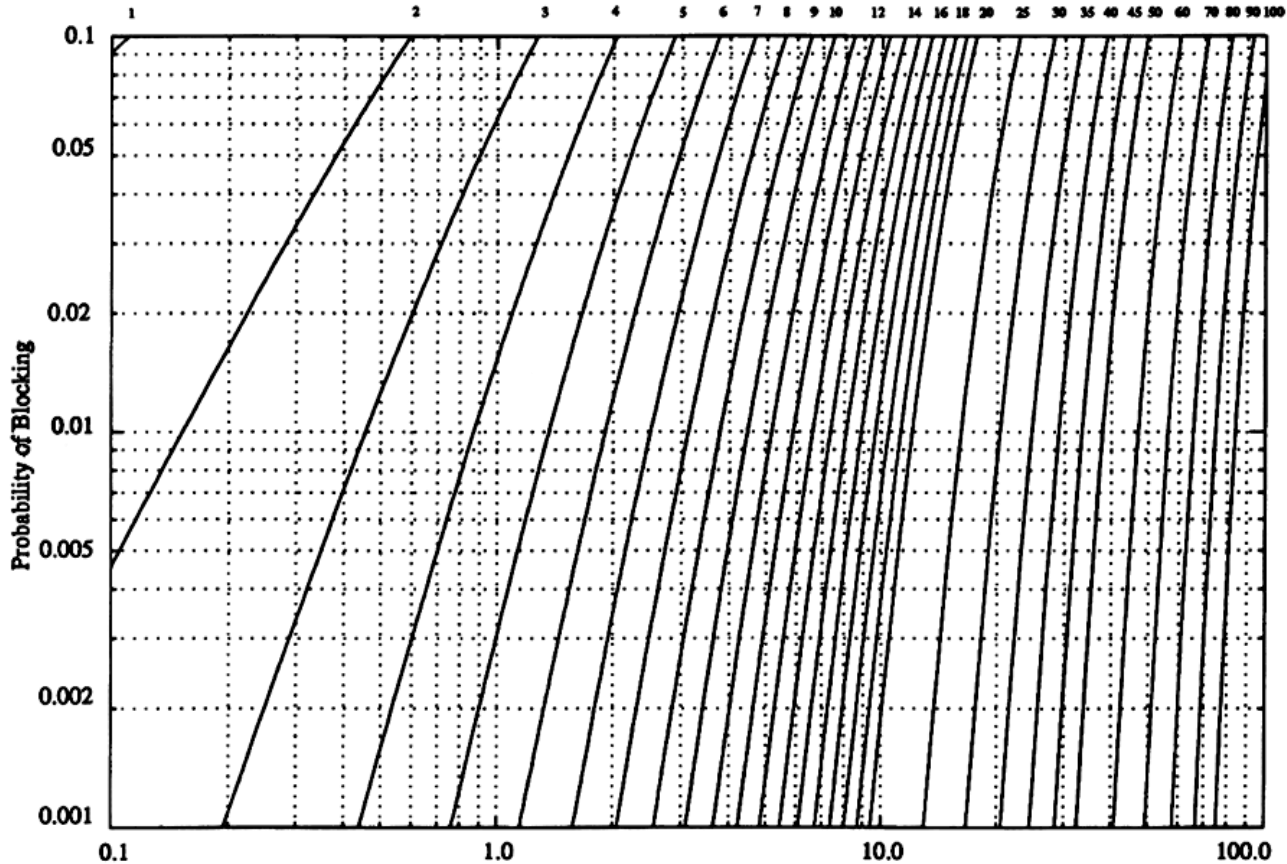
In MATLAB, use `erlangb(m, A)`

*You can download the .m file.*

# Erlang B Formula and Chart

$$P_b = \frac{A^m}{\sum_{i=0}^m \frac{A^i}{i!}}$$

Number of Trunked Channel (m) 



Traffic Intensity in Erlangs (A)

# Example 1

$$P_b \leq 0.005$$

- How many users can be supported for 0.5% blocking probability for the following number of trunked channels in a blocked calls cleared system?
  - (a) 5  $m=5$
  - (b) 10  $m=10$
- Assume each user generates  $A_u = 0.1$  Erlangs of traffic.

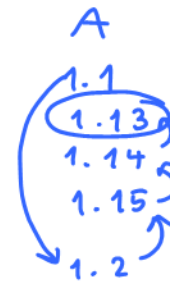
For example,  $\uparrow$

$$A_u = \lambda_u \times \frac{1}{\mu}$$

6 times/day  
average 24 min }  $\Rightarrow \frac{6 \times 24}{24 \times 60} = \frac{1}{10}$

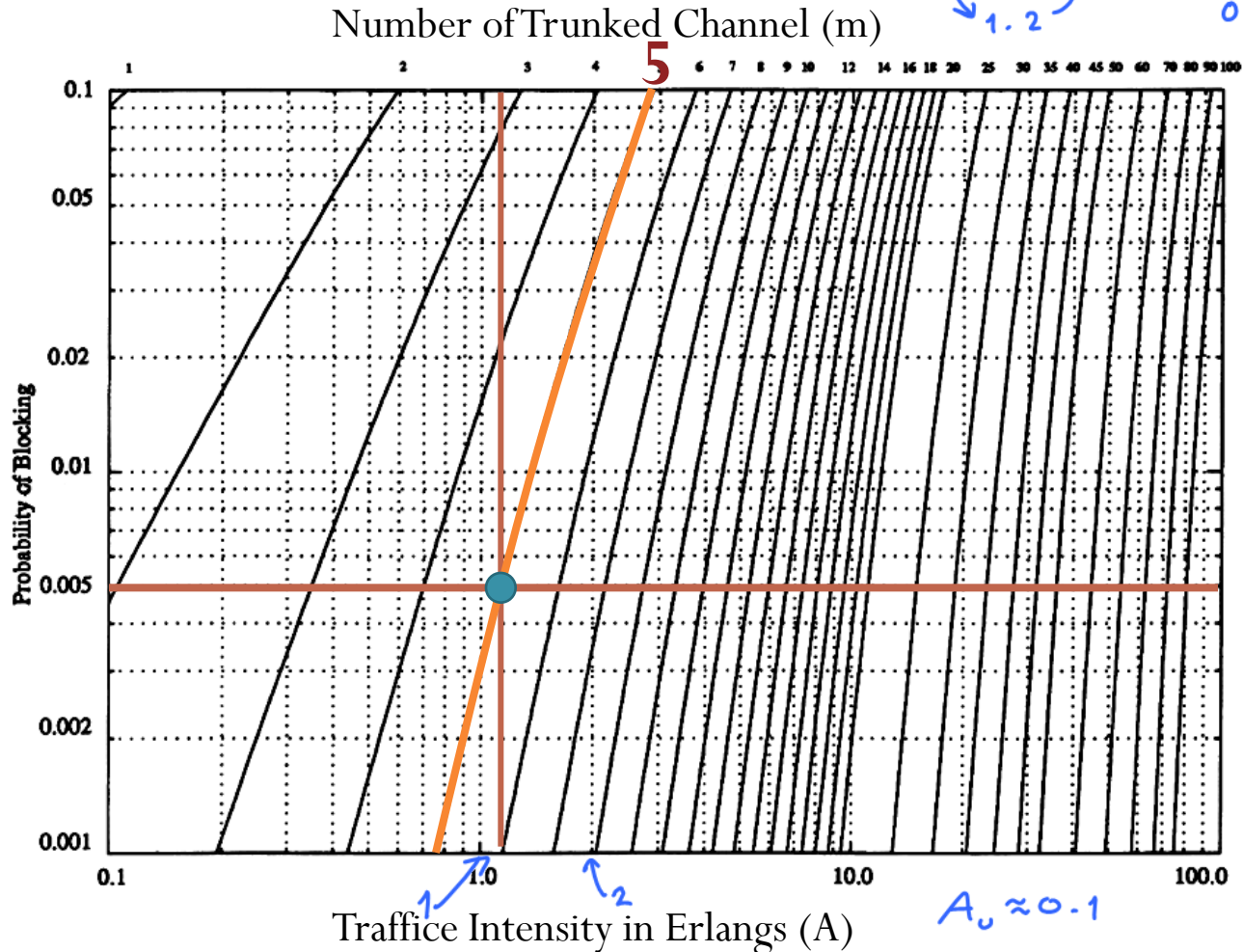
MATLAB  
 $m=5 \rightarrow A=1.13 \rightarrow 11$  users

$m=5$



$P_b$
0.0045
0.0050
0.0051
0.0053
0.0063

# Example 1a



Traffic Intensity in Erlangs (A)

$A_u \approx 0.1$

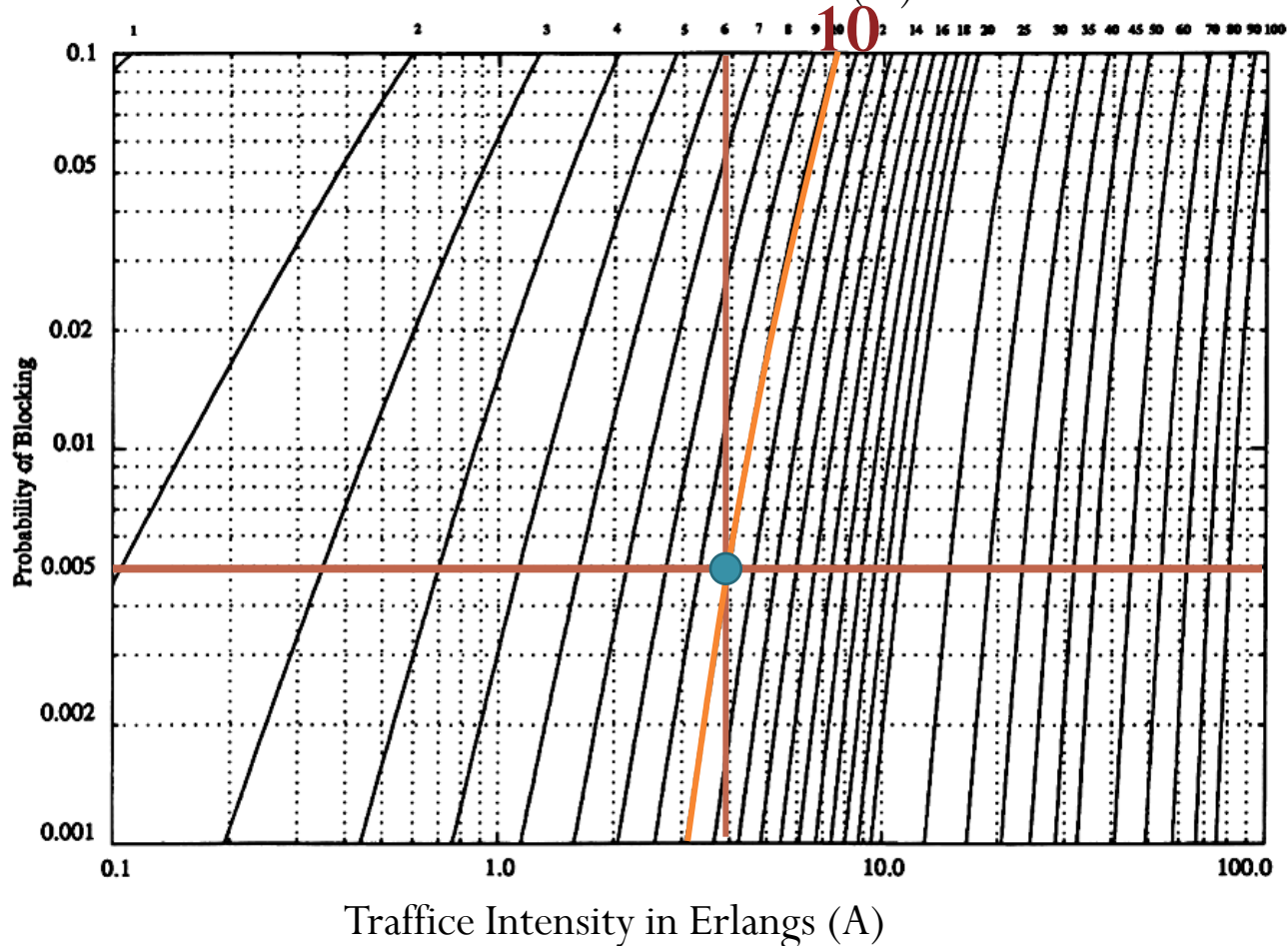
$A = n A_u$

$A \approx 1 \Rightarrow n \approx 10$  users

# Example 1b

$m = 10 \rightarrow A = 3.96 \rightarrow 39 \text{ users}$

Number of Trunked Channel (m)



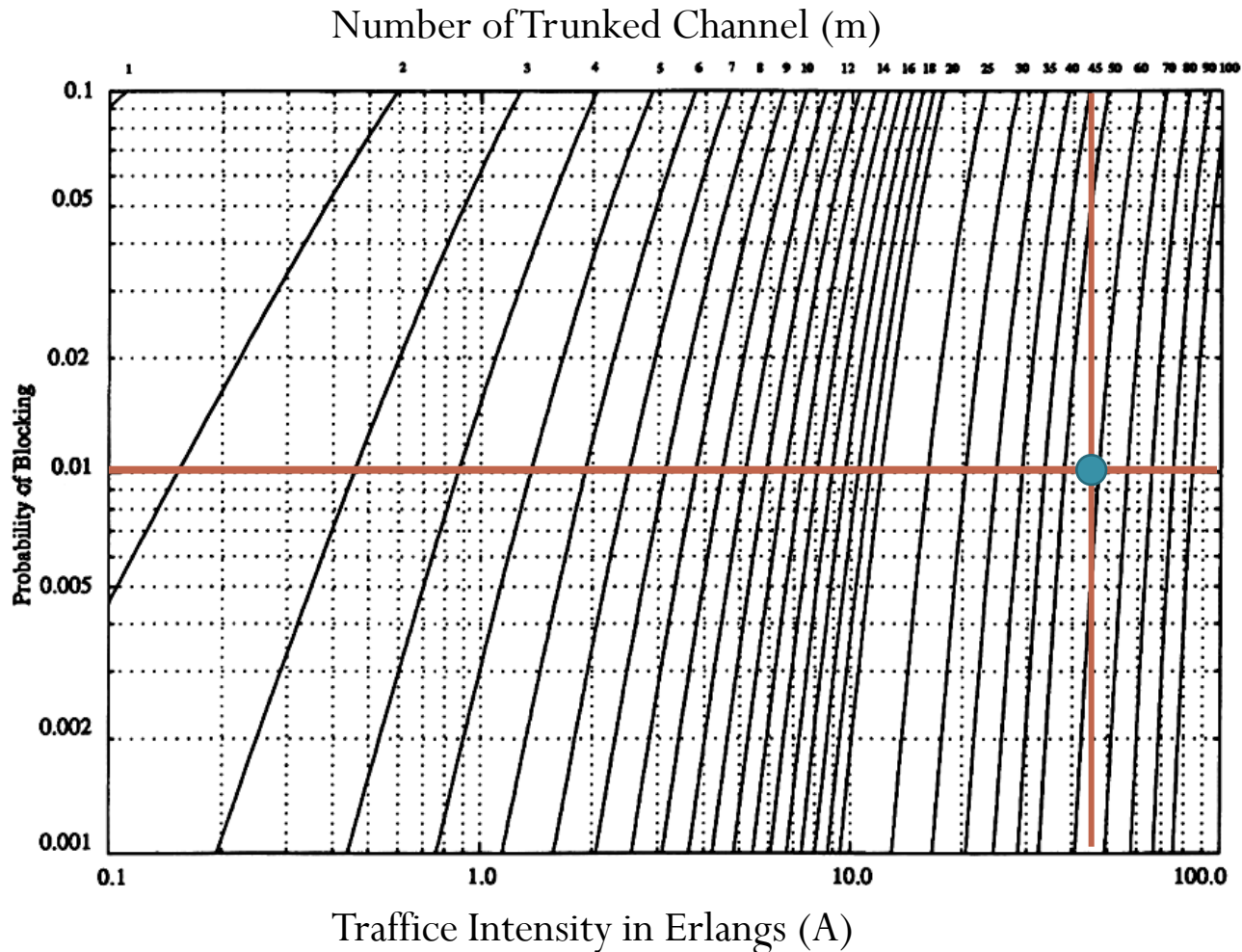
# Example 2.1

- Consider a cellular system in which
  - an average call lasts two minutes  $H = \frac{1}{\mu} = 2 \text{ mins}$
  - the probability of blocking is to be no more than 1%.  $P_b \leq 0.01$
- If there are a total of 395 traffic channels for a seven-cell reuse system, there will be about 57 traffic channels per cell.  $\frac{395}{7} \rightarrow$   $N=7$
- From the Erlang B formula, can handle 44.2 Erlangs or **1326 calls per hour**.

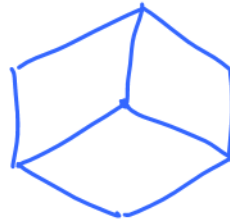
$$A = \lambda \times \frac{1}{\mu}$$
$$44.2 = \lambda \times \frac{2 \text{ mins}}{\text{call}}$$
$$\lambda = \frac{44.2}{2} \frac{\text{calls}}{\text{min}} = 22.1 \frac{\text{calls}}{\text{min}}$$

x60

# Example 2.1: Erlang B



## Example 2.2

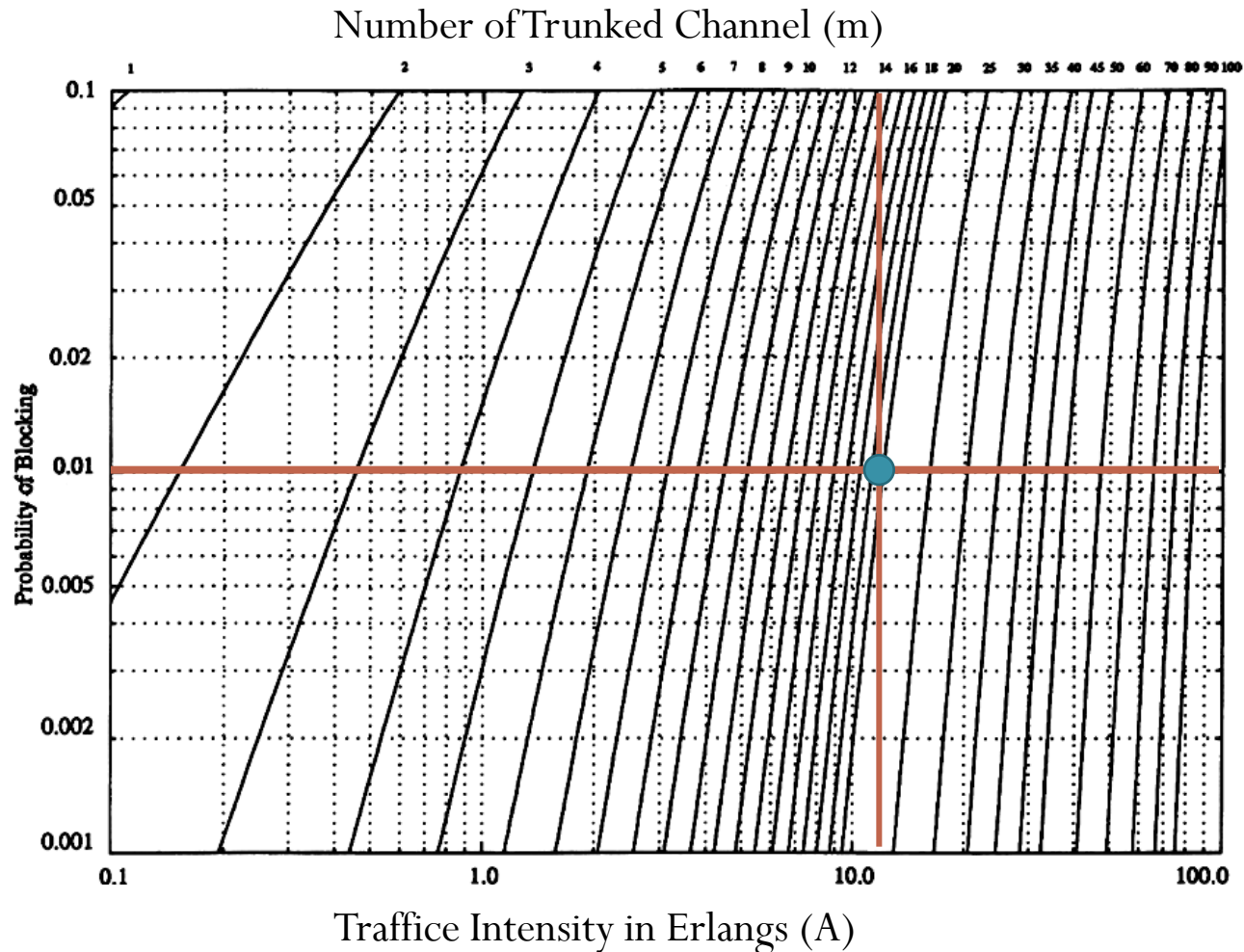


- Now employing **120° sectoring**, there are only 19 channels per sector (57/3 ~~antennas~~ <sup>channel</sup>).
- For the same probability of blocking and average call length, each sector can handle 11.2 Erlangs or 336 calls per hour.  $P_b \leq 0.01$
- Since each cell consists of three sectors, this provides a cell capacity of  $3 \times 336 = 1008$  calls per hour, which amounts to a 24% decrease when compared to the unsectorized case.
- Thus, sectoring decreases the **trunking efficiency** while improving the SIR for each user in the system.

worse Erlang  
but  
better SIR.



# Example 2.2: Erlang B



# Erlang B Trunking Efficiency

**Table 3.4** Capacity of an Erlang B System

Number of Channels $m$	Capacity (Erlangs) for GOS $r_b$			
	1% = 0.01	= 0.005	= 0.002	0.1% = 0.001
2	0.153	0.105	0.065	0.046
4	0.869	0.701	0.535	0.439
5	1.36	1.13	0.900	0.762
10	4.46	3.96	3.43	3.09
20	12.0	11.1	10.1	9.41
24	15.3	14.2	13.0	12.2
40	29.0	27.3	25.7	24.5
70	56.1	53.7	51.0	49.2
100	84.1	80.9	77.4	75.2

m → (points to m=2)  
x 2 (arrow from 10 to 20)  
x > 2 (arrow from 10 to 20)  
A (circles the right side of the table)

# Summary of Chapter 2: Big Picture

$S$  = total # available duplex radio channels for the system

Frequency reuse with **cluster size  $N$**

Path loss exponent

“Capacity”

$$C = \frac{A_{\text{total}}}{A_{\text{cell}}} \times \frac{S}{N} \quad \longleftrightarrow \quad \frac{S}{I} \approx \frac{kR^{-\gamma}}{K \times (kD^{-\gamma})} = \frac{1}{K} \left( \frac{D}{R} \right)^\gamma = \frac{1}{K} \left( \sqrt{3N} \right)^\gamma$$

Tradeoff

$m$  = # channels allocated to each cell.

- Omni-directional:  $K = 6$
- 120° Sectoring:  $K = 2$
- 60° Sectoring:  $K = 1$

Trunking

$m$  = # trunked channels

$\lambda$  = Average # call attempts/requests per unit time

Call blocking probability

$$P_b = \frac{\frac{A^m}{m!}}{\sum_{i=0}^m \frac{A^i}{i!}}$$

$A$  = **traffic intensity** or load [Erlangs] =  $\frac{\lambda}{\mu}$

Erlang-B formula

$\frac{1}{\mu} = H$  = Average call length

# Example 3 (1)

- System Design
- 20 MHz of total spectrum.
- Each simplex channel has 25 kHz RF bandwidth.
- The number of duplex channels:

$$S = \frac{20 \times 10^6}{2 \times 25 \times 10^3} = 400 \text{ channels}$$

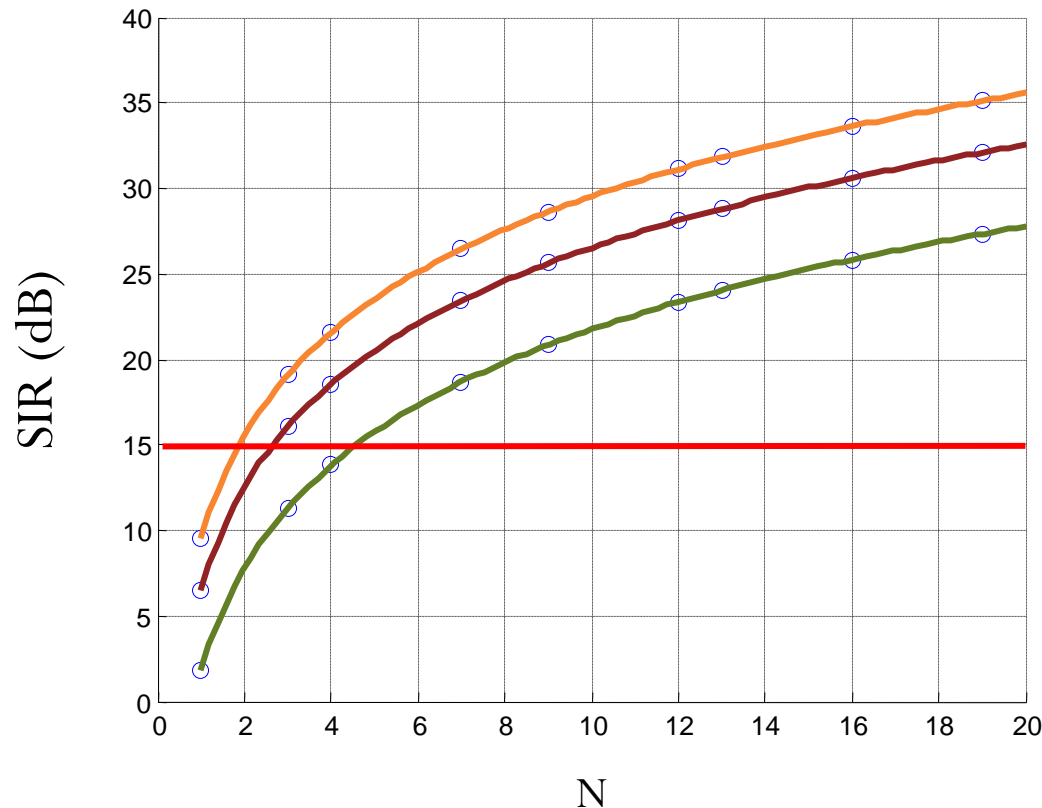
- $\gamma = 4$
- Design requirements:
  - $SIR \geq 15 \text{ dB}$
  - $P_b \leq 5\%$

# Example 3 (2)

- SIR ≥ 15 dB

$$\frac{S}{I} \approx \frac{1}{K} \left( \sqrt{3N} \right)^\gamma$$

```
clear all; close all;
y = 4;
figure; grid on; hold on;
for K = [1,2,6]
    N = [1, 3, 4, 7, 9, 12, 13, 16, 19];
    SIR = 10*log10(1/K*((sqrt(3*N)).^y))
    plot(N,SIR,'o')
N = linspace(1,20,100);
    SIR = 10*log10(1/K*((sqrt(3*N)).^y));
    plot(N,SIR)
end
```



60° sectoring

K = 1 → N = 3

120°

K = 2 → N = 3

Omni

K = 6 → N = 7

# Example 3 (3)

	Omnidirectional	Sectoring (120°)	Sectoring (60°)
K	6	2	1
N	7	3	3
SIR [dB]	18.7	16.1	19.1
$S=400$ #channels/cell	$400/7 = 57$	$400/3 = 133$	$400/3 = 133$
#sectors	1	3	6
$m=$ #channels/sector	57	$133/3 = 44$	$133/6 = 22$
A [Erlangs]/sector	51.55	38.56	17.13
A [Erlangs]/cell	51.55	$38.56 \times 3 = 115.68$	$17.13 \times 6 = 102.78$
#users/cell	18558	41645	37001

Assume that each user makes 2 calls/day and 2 min/call on average  $\rightarrow 1/360$  Erlangs.